

What is a graphical toric arrangement?

Arrangements in Pyrénées

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Chromatic polynomial and flow polynomial

Let G be a graph, V the set of vertices, E the set of edges.

A q -coloring is a map $c : V \rightarrow \mathbb{Z}/q$. It is **proper** if $c(i) \neq c(j) \forall (i, j) \in E$.

FACT: the number $\chi_G(q)$ of proper q -colorings of G is a polynomial in q , which we call the **chromatic polynomial** of G .

FACT (**deletion-contraction**): $\chi_G(q) = \chi_{G \setminus e}(q) - \chi_{G/e}(q)$.

Given an orientation of G , a q -flow is a map $f : E \rightarrow \mathbb{Z}/q$ such that

$$\text{for every } v \in V, \quad \sum_{s(e)=v} f(e) = \sum_{t(e)=v} f(e).$$

It is **nowhere zero** if $f(e) \neq 0 \forall e \in E$.

FACT: the number $\chi_G^*(q)$ of nowhere zero q -flows of G is a polynomial in q , which we call the **flow polynomial** of G .

Again, we have deletion-contraction.

What is most general deletion-contraction invariant of a graph?

The **Tutte polynomial** $T_G(x, y)$.

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The characteristic polynomial

Let \mathcal{H} be a hyperplane arrangement in \mathbb{K}^n .

Its intersection lattice $L(\mathcal{H})$ is a ranked poset, hence it has a Moebius function μ and a **characteristic polynomial**

$$\chi_{\mathcal{H}}(q) = \sum_{L \in L(\mathcal{H})} \mu(L, \hat{0}) q^{\dim(L)}.$$

The polynomial $\chi_{\mathcal{H}}(q)$ contains much information on the complement of the arrangement \mathcal{H} , such as:

- the number of chambers and bounded chambers (if $\mathbb{K} = \mathbb{R}$);
- the Betti numbers of the complement (if $\mathbb{K} = \mathbb{C}$);
- the number of points in the complement (if $\mathbb{K} = \mathbb{F}_q$).

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Graphical hyperplane arrangement

Let $G = (V, E)$ be a graph and $n = |V|$.

To G we associate a **graphical hyperplane arrangement** $\mathcal{H}(G)$ by taking, for every $(i, j) \in E$, the hyperplane of equation $e_i = e_j$ in \mathbb{K}^n .

FACT: $\chi_{\mathcal{H}(G)} = \chi_G$.

We can also define a dual arrangement $\mathcal{H}^*(G)$ such that $\chi_{\mathcal{H}^*(G)} = \chi_G^*$.
Moreover, if G is planar, it exists G^* such that $\mathcal{H}^*(G) = \mathcal{H}(G^*)$.

In recent years, toric arrangements have been studied intensively.
So the question arises naturally:

Problem

What is a graphical toric arrangement?

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Toric arrangements

A list X of elements of a finitely generated abelian group Γ defines a **(generalized) toric arrangement** \mathcal{T}_X ; this is a family of subgroups in the abelian compact Lie group $\text{Hom}(\Gamma, \mathbb{S}^1)$, composed by a subgroup H_e for every $e \in X$ (having codimension 0 if e is torsion and 1 otherwise).

Example: $\Gamma = \mathbb{Z}^2$, $T = \text{Hom}(\Gamma, \mathbb{S}^1) = (\mathbb{S}^1)^2$.

The list $X = \{(2, 0), (0, 3), (-1, 1)\} \subset \mathbb{Z}^2$ defines the toric arrangement \mathcal{T}_X given T by the equations: $t^2 = 1$, $s^3 = 1$, $t^{-1}s = 1$.

The hyperplane arrangement \mathcal{H}_X corresponding to the same list X is given in \mathbb{R}^2 by equations: $2x = 0$, $3y = 0$, $-x + y = 0$.

If we replace $(0, 3)$ by $(0, 5)$, we get the same \mathcal{H}_X , but a different \mathcal{T}_X .

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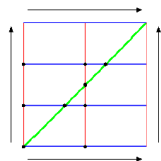
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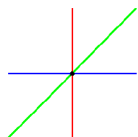
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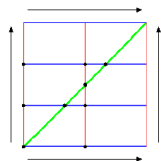
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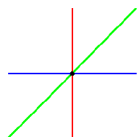
Example: $\Gamma = \mathbb{Z}^2$, $T = \text{Hom}(\Gamma, \mathbb{S}^1) = (\mathbb{S}^1)^2$.

The list $X = \{(2, 0), (0, 3), (-1, 1)\} \subset \mathbb{Z}^2$ defines the toric arrangement \mathcal{T}_X given T by the equations: $t^2 = 1$, $s^3 = 1$, $t^{-1}s = 1$.

The hyperplane arrangement \mathcal{H}_X corresponding to the same list X is given in \mathbb{R}^2 by equations: $2x = 0$, $3y = 0$, $-x + y = 0$.



toric arrangement



hyperplane arrangement

If we replace $(0, 3)$ by $(0, 5)$, we get the same \mathcal{H}_X , but a different \mathcal{T}_X .

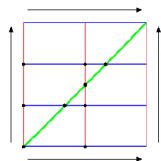
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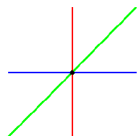
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Matroids and arithmetic matroids

The characteristic polynomial $\chi_{\mathcal{H}}$ of a **hyperplane** arrangement \mathcal{H} only depends on the "linear algebra" of the linear forms defining \mathcal{H} .

This linear algebra is captured by the combinatorial notion of **matroid**.

To every matroid one associates a **Tutte polynomial**, which can be specialized to $\chi_{\mathcal{H}}$: $T(x, y) \doteq \sum_{A \subseteq X} (x-1)^{rk(X)-rk(A)} (y-1)^{|A|-rk(A)}$.

The characteristic polynomial $\chi_{\mathcal{T}}$ of a **toric** arrangement \mathcal{T} depends not only on the linear algebra of the characters defining \mathcal{T} , but also on their arithmetics. We tried to capture the linear algebra AND the arithmetics by the notion of **arithmetic matroid**.

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where in the realizable case $m(A) = |(\Gamma/\langle A \rangle)_{tors}|$.

Has this applications to graph theory?

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Labelled graphs

Graph $\mathcal{G} := (V, E)$ with a map $\ell : E \mapsto \mathbb{Z}_{>0}$ and a partition $E = R \sqcup D$.

For example, let \mathcal{G} , where $\mathcal{G} := (V, E)$, $V := \{v_1, v_2, v_3, v_4\}$,

$R := \{\{v_1, v_2\}, \{v_2, v_3\}, \{v_2, v_4\}\}$ the *regular edges*,

$D := \{\{v_3, v_4\}\}$ the *dotted edges*;

let $\ell(\{v_1, v_2\}) = 1$, $\ell(\{v_2, v_3\}) = 2$, $\ell(\{v_2, v_4\}) = 3$,

$\ell(\{v_3, v_4\}) = 6$ be the *labels*.

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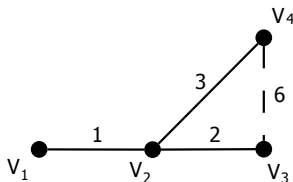
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Oriented labelled graphs

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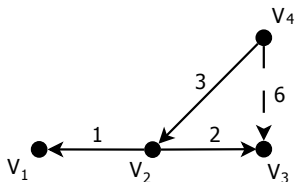
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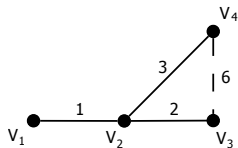
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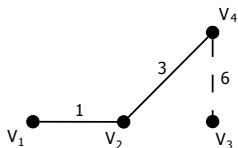
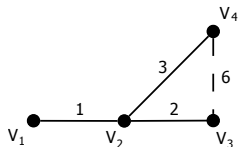
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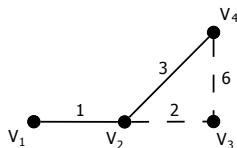
Deletion and contraction



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Deletion of $\{v_2, v_3\}$.



Contraction of $\{v_2, v_3\}$.

Arithmetic colorings

For our results we will consider only positive integers q such that $L(\mathcal{G}) \doteq \text{LCM } \ell(e)$, $e \in E$ divides q . We will call such an integer **admissible**.

A (proper) **arithmetic q -coloring** of a labelled graph \mathcal{G} is a map $c : V \rightarrow \mathbb{Z}/q\mathbb{Z}$ such that:

- (1) if $e := \{u, v\} \in R$, then $\ell(e) \cdot c(u) \neq \ell(e) \cdot c(v)$;
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The **arithmetic chromatic polynomial** $\chi_{\mathcal{G}}(q)$ of \mathcal{G} is defined as the number of (proper) arithmetic q -colorings of \mathcal{G} .

When $D = \emptyset$ and $\ell \equiv 1$ we get the classical chromatic polynomial.

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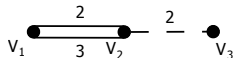
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We have:

$$2c(v_1) \neq 2c(v_2), \quad 3c(v_1) \neq 3c(v_2)$$

$$2c(v_2) = 2c(v_3).$$

We can color v_1 in q ways, then v_2 in $q - 3 - 2 + 1$ ways, then v_3 in 2 ways, so $\chi_{\mathcal{G}}(q) = 2q(q - 4) = 2q^2 - 8q$.



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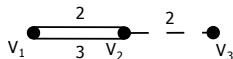
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Arithmetic flows

Given an admissible q , a (nowhere zero) **arithmetic q -flow** on an oriented labelled graph \mathcal{G}_θ is a map $f : E_\theta \rightarrow (\mathbb{Z}/q\mathbb{Z})$ such that:

$$(1) \forall v \in V, \sum_{s(e)=v} \ell(e)f(e) = \sum_{t(e)=v} \ell(e)f(e) \in \mathbb{Z}/q\mathbb{Z}$$

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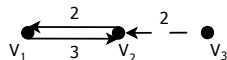
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The equation $2x - 3y = 0$ has q solutions, but 4 of them are not nowhere zero.



We have that $\chi_{\mathcal{G}}^*(q) = 2(q - 4) = 2q - 8$.

Arithmetic flows

Given an admissible q , a (nowhere zero) **arithmetic q -flow** on an oriented labelled graph \mathcal{G}_θ is a map $f : E_\theta \rightarrow (\mathbb{Z}/q\mathbb{Z})$ such that:

$$(1) \forall v \in V, \sum_{s(e)=v} \ell(e)f(e) = \sum_{t(e)=v} \ell(e)f(e) \in \mathbb{Z}/q\mathbb{Z}$$

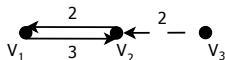
(2) for all $e \in R_\theta$, $f(e) \neq \bar{0} \in \mathbb{Z}/q\mathbb{Z}$.

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Graphical toric arrangements

We associate to \mathcal{G} a list of elements of a group in the following way.

To each edge $e = (v_i, v_j) \in E_\theta$ we associate the element of \mathbb{Z}^n
 $x_e \doteq (0, \dots, 0, \ell(e), 0, \dots, 0, -\ell(e), 0, \dots)$.

Then we look at the image of the list X_R in the group $\Gamma := \mathbb{Z}^n / \langle X_D \rangle$.

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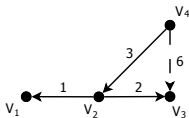
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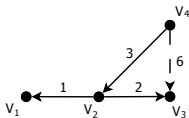
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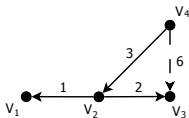
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Main results

Let $\bar{\mathcal{G}} = (\bar{V}, \bar{E})$ be the graph obtained from $\mathcal{G} = (V, E = R \cup D)$ by (classically) contracting the edges in D . Let q be an admissible integer.

Theorem (M.- D'Adderio)

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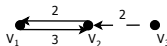
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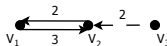
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Arithmetic quasipolynomials

And if q is not ammissible?

In general, the number of arithmetic colorings is a **quasipolynomial** in q .
When $q \equiv 0 \pmod{L(\mathcal{G})}$, it coincides with the arithmetic chromatic polynomial.

When q is coprime with $L(\mathcal{G})$, it coincides with the classical chromatic polynomial.

The same for the number of arithmetic flows.

Together with Petter Brändén, we are now studying an **arithmetic Tutte quasipolynomial** that specializes to the quasipolynomials above, and "interpolates" between the classical and the arithmetic Tutte polynomial. This is defined for any list X in a finitely generated abelian group Γ , in the more general language of "multivariate Tutte polynomials":

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